

# Computing modular polynomials in dimension 2

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# Computing modular polynomials

- 1 Dimension 1 : elliptic curves
- 2 Dimension 2 : abelian surfaces
  - Computation of the modular polynomials
  - Smaller invariants
- 3 Real Multiplication : cyclic isogenies

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# Motivation

An **isogeny** between two elliptic curves  $E_1$  and  $E_2$  is a surjective map with finite kernel. The **degree** of the isogeny is the cardinality of the kernel.

Many applications :

- Theory ;
- Cryptography : transfert the DLP ;
- SEA algorithm ;
- Class polynomials ;
- Graph of isogenies.

# Motivation

For cryptography, we work over finite fields.

Here, we work on  $\mathbb{C}$ .

- The theory is “easy” on  $\mathbb{C}$  ;
- Numerical computation ;
- The modular polynomials can be reduced modulo  $p$ .

# Complex elliptic curves

Let  $\mathcal{H}_1 = \{a + \imath b : b > 0\} \subset \mathbb{C}$  be the **Poincaré half plane**.

## Proposition

Let  $E/\mathbb{C}$  be an elliptic curve.

Then there exists a lattice

$$\Lambda = \mathbb{Z} + \tau\mathbb{Z}, \quad \text{where } \tau \in \mathcal{H}_1$$

and a complex analytic isomorphism  $E \simeq \mathbb{C}/\Lambda$  of complex Lie groups.

# Isomorphism

**Modular group** :  $SL_2(\mathbb{Z}) = \left\{ \begin{pmatrix} a & b \\ c & d \end{pmatrix} : a, b, c, d \in \mathbb{Z}, ad - bc = 1 \right\}$

**Group action** :

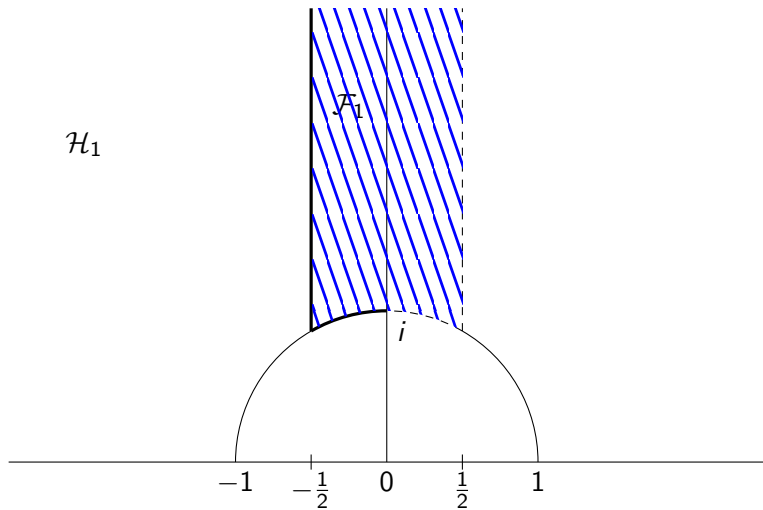
$$\begin{aligned} SL_2(\mathbb{Z}) \times \mathcal{H}_1 &\longrightarrow \mathcal{H}_1 \\ \begin{pmatrix} a & b \\ c & d \end{pmatrix} \tau &= \frac{a\tau + b}{c\tau + d} \end{aligned}$$

## Proposition

Two elliptic curves  $E_1$  and  $E_2$  over  $\mathbb{C}$  corresponding to the lattices  $\Lambda_1 = \mathbb{Z} + \tau_1\mathbb{Z}$  and  $\Lambda_2 = \mathbb{Z} + \tau_2\mathbb{Z}$  are **isomorphic** if and only if there exists  $\gamma \in SL_2(\mathbb{Z})$  such that  $\tau_2 = \gamma\tau_1$ .

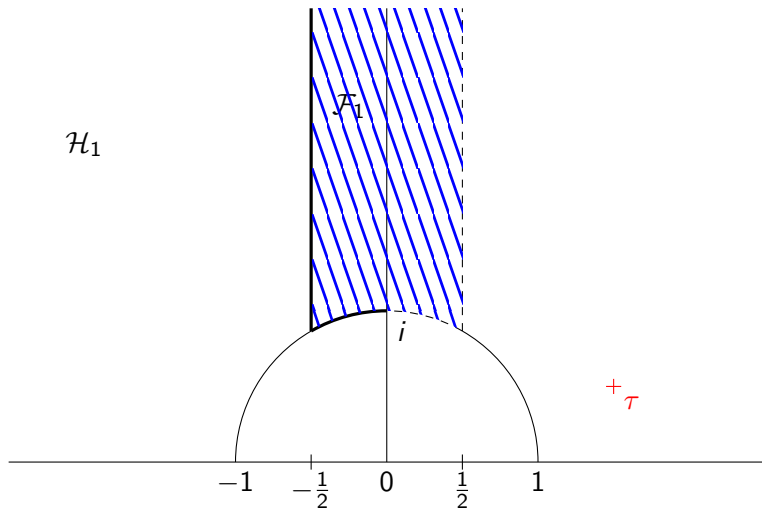
$\implies$  change of basis of the lattice.

# Fundamental domain $\mathcal{F}_1$

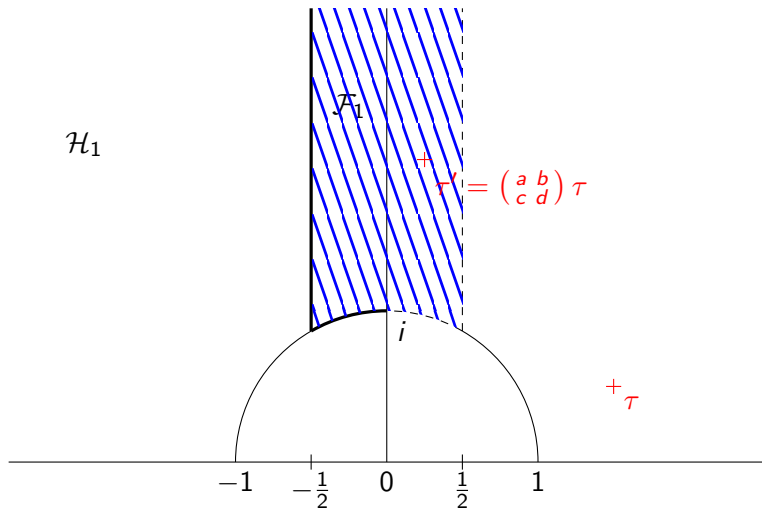




# Fundamental domain $\mathcal{F}_1$



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# Modular function

Let  $p$  be a prime and

$$\Gamma_0(p) = \left\{ \begin{pmatrix} a & b \\ c & d \end{pmatrix} \in \mathrm{SL}_2(\mathbb{Z}) : c \equiv 0 \pmod{p} \right\}.$$

## Definition

Let  $\Gamma \subset \mathrm{SL}_2(\mathbb{Z})$  be a congruence subgroup of finite index.

$f : \mathcal{H}_1 \rightarrow \mathbb{C}$  is a **modular function** for  $\Gamma$  if

- 1  $f$  is meromorphic on  $\mathcal{H}_1$  (and on the cusps);
- 2 for all  $\gamma \in \Gamma$  and  $\tau \in \mathcal{H}_1$ ,  $f(\gamma\tau) = f(\tau)$ .

## Example

- $j(\tau)$  is a modular function for  $\mathrm{SL}_2(\mathbb{Z})$ ;
- $j_p(\tau) := j(p\tau)$  is a modular function for  $\Gamma_0(p)$ .

# Modular function

Let  $\Gamma \subset \mathrm{SL}_2(\mathbb{Z})$  be a congruence subgroup of finite index. Denote by  $\mathbb{C}_\Gamma$  the field of modular functions for  $\Gamma$ . Then

## Theorem

- $\mathbb{C}_{\mathrm{SL}_2(\mathbb{Z})} = \mathbb{C}(j)$ ;
- $\mathbb{C}_{\Gamma_0(p)} = \mathbb{C}(j, j_p)$ .

# Isogeny

We are interested in the isogenies of degree  $p$ .

Let  $C_p$  be a set of representatives of  $SL_2(\mathbb{Z})/\Gamma_0(p)$ .

The isogenous points of degree  $p$  are :  $p\gamma\tau$ ,  $\gamma \in C_p$ .

## Theorem

*The field extension  $\mathbb{C}_{\Gamma_0(p)}/\mathbb{C}_{SL_2(\mathbb{Z})} = \mathbb{C}(j, j_p)/\mathbb{C}(j)$  is algebraic of degree  $[SL_2(\mathbb{Z}) : \Gamma_0(p)] = p + 1$ .*

Conjugate functions of  $j_p$  in  $\mathbb{C}_{\Gamma_0(p)}/\mathbb{C}_{SL_2(\mathbb{Z})}$  :  $j_p^\gamma(\tau) := j(p\gamma\tau)$ ,  $\gamma \in C_p$ .

# Modular polynomial

The **classical modular polynomial** of index  $p$  is the polynomial  $\Phi_p$  that parameterizes isomorphism classes of elliptic curves together with an isogeny of degree  $p$  :

$$\Phi_p(X, j(E)) = \prod_{E' \text{ } p\text{-isogenous to } E} (X - j(E')).$$

It is also the minimal polynomial of  $j_p$  for the extension  $\mathbb{C}_{\Gamma_0(p)}/\mathbb{C}_{\text{SL}_2(\mathbb{Z})}$ , thus

$$\Phi_p(X, j) = \prod_{\gamma \in C_p} (X - j_p^\gamma) \in \mathbb{Z}[X, j].$$

## Algorithm

Computation of the modular polynomials by evaluation/interpolation (Enge 2009).

$$\Phi_p(X, j) = \prod_{\gamma \in \mathbb{C}_p} (X - j^\gamma) = X^{p+1} + \sum_{i=0}^p c_i(j) X^i.$$

- Evaluate :

$$\prod_{\gamma \in \mathbb{C}_p} (X - j(p\gamma\tau)) = X^{p+1} + \sum_{i=0}^p c_i(j(\tau)) X^i;$$

$\Rightarrow$  Evaluate in  $\deg_j(\Phi_p) + 1 = (p + 1) + 1$  values  $\tau$ .

- Interpolate  $c_i$ .

Evaluation of  $j$  in  $\tilde{O}(N)$  at precision  $N$  digits (Dupont 2006);  
Algorithm quasi-linear :  $\tilde{O}(p^3)$ .

# Examples

$$p = 2 \quad \Phi_2(X, Y) = X^3 + (-Y^2 + 1488Y - 162000)X^2 + (1488Y^2 + 40773375Y + 8748000000)X + (Y^3 - 162000Y^2 + 8748000000Y - 157464000000000)$$

$$p = 3 \quad \Phi_3(X, Y) = X^4 + (-Y^3 + 2232Y^2 - 1069956Y + 36864000)X^3 + (2232Y^3 + 2587918086Y^2 + 8900222976000Y + 452984832000000)X^2 + (-1069956Y^3 + 8900222976000Y^2 - 770845966336000000Y + 185542587187200000000)X + (Y^4 + 36864000Y^3 + 452984832000000Y^2 + 1855425871872000000000Y)$$

$$p = 5 \quad \Phi_5(X, Y) = X^6 + (-Y^5 + 3720Y^4 - 4550940Y^3 + 2028551200Y^2 - 246683410950Y + 1963211489280)X^5 + (3720Y^5 + 1665999364600Y^4 + 107878928185336800Y^3 + 383083609779811215375Y^2 + 128541798906828816384000Y + 1284733132841424456253440)X^4 + (-4550940Y^5 + 107878928185336800Y^4 - 441206965512914835246100Y^3 + 26898488858380731577417728000Y^2 - 192457934618928299655108231168000Y + 280244777828439527804321565297868800)X^3 + (2028551200Y^5 + 383083609779811215375Y^4 + 26898488858380731577417728000Y^3 + 5110941777552418083110765199360000Y^2 + 36554736583949629295706472332656640000Y + 6692500042627997708487149415015068467200)X^2 + (-246683410950Y^5 + 128541798906828816384000Y^4 - 192457934618928299655108231168000Y^3 + 36554736583949629295706472332656640000Y^2 - 264073457076620596259715790247978782949376Y + 53274330803424425450420160273356509151232000)X + (Y^6 + 1963211489280Y^5 + 1284733132841424456253440Y^4 + 280244777828439527804321565297868800Y^3 + 6692500042627997708487149415015068467200Y^2 + 53274330803424425450420160273356509151232000Y + 141359947154721358697753474691071362751004672000)$$

Other invariants : Schläfli, Weber, theta functions.

Schläfli 1870,  $p = 5$  :  $x^6 - x^5y^5 + 4xy + y^6$ .



# Computing modular polynomials

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# Motivation

Dimension 2 : principally polarized abelian surfaces (ppas)  $\implies$  Jacobian of hyperelliptic curves of genus 2 (or product of elliptic curves);

Cryptography : competitive with elliptic curves ;

$\implies$  we want to do the same thing !

## Siegel space

**Siegel upper half-space**  $\mathcal{H}_2$  the set of  $2 \times 2$  symmetric matrices over  $\mathbb{C}$  with positive definite imaginary part.

Ppas on  $\mathbb{C}$  :  $A \simeq \mathbb{C}^2/\Lambda$  where  $\Lambda = \mathbb{Z}^2 + \Omega\mathbb{Z}^2$ , with  $\Omega \in \mathcal{H}_2$  (**period matrix**).

Let  $J = \begin{pmatrix} 0 & Id_2 \\ -Id_2 & 0 \end{pmatrix}$ . **Symplectic group** :

$$\mathrm{Sp}_4(\mathbb{Z}) = \{\gamma \in \mathrm{GL}_4(\mathbb{Z}) : {}^t\gamma J \gamma = J\}.$$

**Group action** :  $\begin{pmatrix} A & B \\ C & D \end{pmatrix} \Omega = (A\Omega + B)(C\Omega + D)^{-1}$ .

We have a fundamental domain  $\mathcal{F}_2$ .

**Siegel modular threefold** :  $\mathcal{H}_2/\mathrm{Sp}_4(\mathbb{Z})$ .

# Modular forms and functions

Let  $\Gamma$  be a subgroup of finite index of  $\mathrm{Sp}_4(\mathbb{Z})$  and  $k \in \mathbb{Z}$ .

## Definition

A **Siegel modular form** of weight  $k$  for  $\Gamma$  is a function  $f : \mathcal{H}_2 \rightarrow \mathbb{C}$  such that :

- 1  $f$  is holomorphic on  $\mathcal{H}_2$  ;
- 2  $f(\gamma\Omega) = \det(C\Omega + D)^k f(\Omega)$ ,  $\forall \gamma = \begin{pmatrix} A & B \\ C & D \end{pmatrix} \in \Gamma$  and  $\Omega \in \mathcal{H}_2$ .

## Definition

**Siegel modular function** for  $\Gamma$  :  $f = \frac{f_1}{f_2}$  quotient of Siegel modular forms of same weight. Thus,  $f(\gamma\Omega) = f(\Omega)$ .

# Theta functions

**Theta functions** (of characteristic  $\frac{1}{2}$ ) :

Define for  $a = (a_0, a_1)$  and  $b = (b_0, b_1)$  in  $\{0, 1\}^2$  :

$$\theta_{b_0+2b_1+4a_0+8a_1}(\Omega) = \sum_{n \in \mathbb{Z}^2} \exp(i\pi {}^t(n + \frac{a}{2})\Omega(n + \frac{a}{2}) + i\pi {}^t(n + \frac{a}{2})b)$$

- 16 theta functions ;
- 6 are identically zero ;
- $\mathcal{P} = \{0, 1, 2, 3, 4, 6, 8, 9, 12, 15\}$  ;
- $\theta_i^2 =$  Siegel modular form of weight 1 for  $\Gamma(2, 4)$ .

$$\Gamma(2, 4) = \left\{ \begin{pmatrix} A & B \\ C & D \end{pmatrix} \in \mathrm{Sp}_4(\mathbb{Z}) : \begin{pmatrix} A & B \\ C & D \end{pmatrix} \equiv \mathrm{Id}_4 \pmod{2}, B_0 \equiv C_0 \equiv 0 \pmod{4} \right\}.$$

# Theta functions

Let

$$h_{10} = \prod_{i \in \mathcal{P}} \theta_i^2,$$

$$h_4 = \sum_{i \in \mathcal{P}} \theta_i^8,$$

$$h_6 = \sum_{60 \text{ triples } (i,j,k) \in \mathcal{P}^3} \pm (\theta_i \theta_j \theta_k)^4,$$

$$h_{12} = \sum_{15 \text{ tuples } (i,j,k,l,m,n) \in \mathcal{P}^6} (\theta_i \theta_j \theta_k \theta_l \theta_m \theta_n)^4.$$

$\Rightarrow h_i$  is a Siegel modular form of weight  $i$  for the group  $\mathrm{Sp}_4(\mathbb{Z})$ .

# Generalization of the $j$ -invariant

## Definition

We call **Igusa invariants**, or  $j$ -invariants, the Siegel modular functions  $j_1, j_2, j_3$  for  $\mathrm{Sp}_4(\mathbb{Z})$  defined by

$$j_1 := \frac{h_{12}^5}{h_{10}^6}, \quad j_2 := \frac{h_4 h_{12}^3}{h_{10}^4}, \quad j_3 := \frac{(h_{12} h_4 - 2h_6 h_{10}) h_{12}^2}{3h_{10}^4}.$$

## Theorem (Igusa 1962, Spallek 1994)

The field of Siegel modular functions invariant by  $\mathrm{Sp}_4(\mathbb{Z})$  is  $\mathbb{C}(j_1, j_2, j_3)$ .

Generically, two ppas have the same  $j$ -invariants if and only if they are isomorphic.



# Isogeny

The functions

$$j_{\ell,p}(\Omega) := j_{\ell}(p\Omega), \quad \ell = 1, 2, 3,$$

are Siegel modular functions for  $\Gamma_0(p)$ , where

$$\Gamma_0(p) := \left\{ \begin{pmatrix} A & B \\ C & D \end{pmatrix} \in \mathrm{Sp}_4(\mathbb{Z}) : C \equiv 0 \pmod{p} \right\}$$

is of index  $[\mathrm{Sp}_4(\mathbb{Z}) : \Gamma_0(p)] = p^3 + p^2 + p + 1$ .

$p$ pas  $(p, p)$ -isogenous to  $\Omega : p\gamma\Omega$ , where  $\gamma \in C_p = \mathrm{Sp}_4(\mathbb{Z})/\Gamma_0(p)$ .

## Theorem (Bröker-Lauter 2009)

*The field of Siegel modular functions invariant by  $\Gamma_0(p)$  is  $\mathbb{C}(j_1, j_2, j_3, j_{1,p})$ .*

We define

$$j_{\ell,p}^{\gamma}(\Omega) := j_{\ell}(p\gamma\Omega), \quad \ell = 1, 2, 3.$$

## Modular polynomials in dimension 2

$$\Phi_{1,p}(X) = \prod_{\gamma \in C_p} (X - j_{1,p}^\gamma),$$

(minimal polynomial of the extension  $\mathbb{C}(j_1, j_2, j_3, j_{1,p})/\mathbb{C}(j_1, j_2, j_3)$ ),

and for  $\ell = 2, 3$ , 
$$\Psi_{\ell,p}(X) = \sum_{\gamma \in C_p} j_{\ell,p}^\gamma \prod_{\gamma' \in C_p \setminus \{\gamma\}} (X - j_{1,p}^{\gamma'}).$$

Proposition (Bröker-Lauter 2009)

$$\Phi_{1,p}, \Psi_{2,p}, \Psi_{3,p} \in \mathbb{Q}(j_1, j_2, j_3)[X].$$

We have  $j_{\ell,p}^\gamma(\Omega) =$

$$\Psi_{\ell,p}(j_{1,p}^\gamma(\Omega), j_1(\Omega), j_2(\Omega), j_3(\Omega)) / \Phi'_{1,p}(j_{1,p}^\gamma(\Omega), j_1(\Omega), j_2(\Omega), j_3(\Omega)).$$

# Algorithm

How to compute the modular polynomials ?

Evaluation/interpolation :

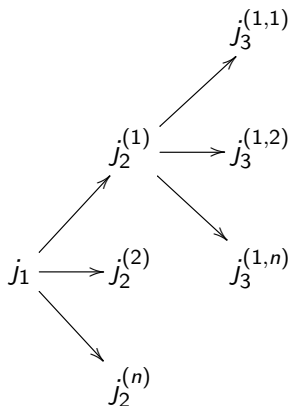
$$\Phi_{1,\rho}(X, j_1(\Omega), j_2(\Omega), j_3(\Omega)) = X^{\rho^3+\rho^2+\rho+1} + \sum_{i=0}^{\rho^3+\rho^2+\rho} c_i(j_1(\Omega), j_2(\Omega), j_3(\Omega)) X^i.$$

where  $c_i \in \mathbb{Q}(j_1, j_2, j_3)$ .

Problem : Interpolation of trivariate rational fractions.

# Interpolation

We can not choose the matrices  $\Omega$  as we want !



# Complexity

Inversion :  $(j_1(\Omega), j_2(\Omega), j_3(\Omega)) \longrightarrow \Omega$  in  $\tilde{O}(N)$  (Dupont 2006).

Fast evaluation of the Igusa invariants :  $\tilde{O}(N)$  (Dupont 2006, Enge–Thomé 2014).

Complexity of the computation of the modular polynomials :  
 $\tilde{O}(d_{j_1} d_{j_2} d_{j_3} p^3 N)$ .

## Streng invariants

The modular polynomials have been computed by Dupont for  $p = 2$  only. For  $p = 3$  they are too big.

### Other invariants?

⇒ Streng 2010 :

$$i_1 := \frac{h_4 h_6}{h_{10}}, \quad i_2 := \frac{h_4^2 h_{12}}{h_{10}^2}, \quad i_3 := \frac{h_4^5}{h_{10}^2}.$$

while Igusa :

$$j_1 := \frac{h_{12}^5}{h_{10}^6}, \quad j_2 := \frac{h_4 h_{12}^3}{h_{10}^4}, \quad j_3 := \frac{(h_{12} h_4 - 2 h_6 h_{10}) h_{12}^2}{3 h_{10}^4}.$$

### Theorem

*The field of Siegel modular functions is  $\mathbb{C}(j_1, j_2, j_3) = \mathbb{C}(i_1, i_2, i_3)$ .*

## Comparison

For  $p = 3$  :  $\Phi_{1,3}(X, f_1, f_2, f_3) = X^{40} + \sum_{i=0}^{39} c_i(f_1, f_2, f_3)X^i$ .

| $i$      | $j_1$    | $i_1$ | $j_2$    | $i_2$ | $j_3$ | $i_3$    |
|----------|----------|-------|----------|-------|-------|----------|
| 0        | 394      | 61    | 288      | 32    | 278   | 32       |
| 1        | 302      | 61    | 286      | 32    | 276   | 31       |
| 2        | 302      | 61    | 286      | 32    | 276   | 31       |
| $\vdots$ | $\vdots$ |       | $\vdots$ |       |       | $\vdots$ |
| 37       | 268      | 41    | 382      | 22    | 253   | 21       |
| 38       | 263      | 36    | 375      | 21    | 248   | 19       |
| 39       | 257      | 31    | 367      | 20    | 243   | 17       |

Memory space :

- $p = 2$  : 2.1 MB against 57 MB.
- $p = 3$  : 890 MB.

# Denominators

- Denominators for **Igusa** invariants for  $p = 2$  :

$$j_1^\alpha D_2(j_1, j_2, j_3)^6 \quad (\alpha \text{ ranging between } 5 \text{ and } 21)$$

- Denominators for **Streng** invariants for  $p = 2$  :

$$i_3^\alpha D_2(i_1, i_2, i_3) \text{ and } i_3^\alpha D_2(i_1, i_2, i_3)^2 \quad (\alpha \text{ varies from } 0 \text{ to } 3)$$

- Denominators for **Streng** invariants for  $p = 3$  :

$$i_3^\alpha D_3(i_1, i_2, i_3)^2 \text{ and } i_3^\alpha D_3(i_1, i_2, i_3)^4 \quad (\alpha \text{ varies from } 0 \text{ to } 4)$$



## Examples

$$D_2(Ig) = 236196j_1^5 + (-972j_2^2 + (5832j_3 + 19245600)j_2 + (-8748j_3^2 - 104976000j_3 + 125971200000))j_1^4 + (j_2^4 + (-12j_3 - 77436)j_2^3 + (54j_3^2 + 870912j_3 - 507384000)j_2^2 + (-108j_3^3 - 3090960j_3^2 + 2099520000j_3)j_2 + (81j_3^4 + 349920j_3^3))j_1^3 + (78j_2^5 + (-1332j_3 + 592272)j_2^4 + (8910j_3^2 - 4743360j_3)j_2^3 + (-29376j_3^3 + 9331200j_3^2)j_2^2 + 47952j_3^4j_2 - 31104j_3^5)j_1^2 + (-159j_2^6 + (1728j_3 - 41472)j_2^5 - 6048j_3^2j_2^4 + 6912j_3^3j_2^3)j_1 + (80j_2^7 - 384j_3j_2^6).$$

$$D_2(Str) = (24576i_3i_1^5 + (96i_2^3 - 4608i_3i_2)i_1^4 + (-6220800i_3i_2 - 12288i_3^2)i_1^3 + (-23328i_2^4 - 48i_3i_2^3 + 1088640i_3i_2^2 + 2304i_3^2i_2 + 24883200i_3^2)i_1^2 + (93312i_3i_2^3 + 419904000i_3i_2^2 - 5909760i_3^2i_2 + (1536i_3^3 - 8398080000i_3^2))i_1 + (1417176i_2^5 - 5832i_3i_2^4 + (6i_2^2 - 94478400i_3)i_2^3 + 287712i_2^2i_2^2 + (-288i_3^3 + 1154736000i_3^2)i_2 + (-248832i_3^3 + 755827200000i_3^2))).$$

$D_3(Str) = 1073741824i^{12}i_2i_3 + 1048576i^{11}i_2i_3 - 100663296i^{10}i_2i_3 -$   
 $805306368i^{12}i_2 + 23653961975736i^{11}i_2i_3 - 16106127236i^{10}i_2i_3 + 391378894848i^{11}i_2 +$   
 $23123460096i^{10}i_2 - 1572864i^{10}i_2i_3 - 222871385088i^{10}i_2i_3 + 150994944i^{10}i_2i_3 +$   
 $1912538199490560i^{10}i_2i_3 + 120795952i^{10}i_2i_3 - 3962711103360i^{10}i_2i_3 + 1885039755264i^{10}i_2i_3 +$   
 $2170276644124486i^{10}i_2i_3 - 134649922477184i^{10}i_2i_3 + 100632960i^{10}i_2i_3 -$   
 $152617815730248744960i^{10}i_2i_3 - 14481997556631212i^{10}i_2i_3 + 212468467875840i^{10}i_2i_3 +$   
 $65691848000i^{10}i_2i_3 + 983040i^{10}i_2i_3 - 1694880688432742i^{10}i_2i_3 - 63031521188160i^{10}i_2i_3 -$   
 $94371840i^{10}i_2i_3 + 352004832446119936i^{10}i_2i_3 - 853811062209536i^{10}i_2i_3 +$   
 $219210284369345614184480i^{10}i_2i_3 - 754974720i^{10}i_2i_3 + 11517392956354127872i^{10}i_2i_3 -$   
 $13232363778146304i^{10}i_2i_3 - 784286613504i^{10}i_2i_3 + 320194663131271231916i^{10}i_2i_3 +$   
 $1917038253511802880i^{10}i_2i_3 + 21388245073920i^{10}i_2i_3 + 19073607346394579179296i^{10}i_2i_3 -$   
 $335544320i^{10}i_2i_3 + 29804560874959987712i^{10}i_2i_3 + 6025767309083052i^{10}i_2i_3 -$   
 $1699781338847946744004608i^{10}i_2i_3 + 104120172682553920i^{10}i_2i_3 + 63490989219840i^{10}i_2i_3 -$   
 $52624982016i^{10}i_2i_3 + 28332151878881801024i^{10}i_2i_3 - 327680i^{10}i_2i_3 + 2205266003248272i^{10}i_2i_3 +$   
 $5051361263616i^{10}i_2i_3 - 1909382002445824718929929i^{10}i_2i_3 - 31457280i^{10}i_2i_3 -$   
 $3282625107925982080i^{10}i_2i_3 - 7784410537473661314263900160i^{10}i_2i_3 +$   
 $810960685032964i^{10}i_2i_3 - 626888273473709361218656i^{10}i_2i_3 + 251658240i^{10}i_2i_3 -$   
 $215987170489694945280i^{10}i_2i_3 + 80465457603017781129543680i^{10}i_2i_3 - 15017529555950548096i^{10}i_2i_3$   
 $10569523360776192i^{10}i_2i_3 - 4638615481599896235073536i^{10}i_2i_3 - 118674948096i^{10}i_2i_3 +$   
 $950395740935763952640i^{10}i_2i_3 + 2941994259840540672i^{10}i_2i_3 + 36187794384906526379541872644$   
 $826872568170644i^{10}i_2i_3 + 155755908742384175609856i^{10}i_2i_3 + 62914560i^{10}i_2i_3 -$   
 $8667079688200173056i^{10}i_2i_3 - 782109458301523924713308956i^{10}i_2i_3 +$   
 $91081207185408i^{10}i_2i_3 + 45662207265162884601085952i^{10}i_2i_3 - 80900057011133778805681037312i^{10}i_2i_3 +$   
 $2870105097548602679040i^{10}i_2i_3 + 896261713716314880i^{10}i_2i_3 + 18809714935249224922i^{10}i_2i_3 +$   
 $4504048233864351947764i^{10}i_2i_3 + 11221327872i^{10}i_2i_3 + 780494067727704797184i^{10}i_2i_3 +$   
 $61440i^{10}i_2i_3 - 273088658240188416i^{10}i_2i_3 - 3987082393876826603145984i^{10}i_2i_3 +$   
 $107088223232i^{10}i_2i_3 + 8443058318853636251648i^{10}i_2i_3 + 4652016478805165040820933736448i^{10}i_2i_3$   
 $58924040i^{10}i_2i_3 - 133879677901996032i^{10}i_2i_3 - 36831815040167945837868423168i^{10}i_2i_3 +$   
 $82081531035648i^{10}i_2i_3 + 2449653429540496482238464i^{10}i_2i_3 - 181127172379191892949216756695$   
 $47185920i^{10}i_2i_3 + 71004250037730410496i^{10}i_2i_3 - 56148085403579770990704795648i^{10}i_2i_3 -$   
 $193303195584228338932224i^{10}i_2i_3 + 90123752946066684928i^{10}i_2i_3 + 8625368264019523372228416$   
 $7400456973692928i^{10}i_2i_3 - 1339816229728639745575936i^{10}i_2i_3 + 88576229376i^{10}i_2i_3 -$   
 $14192373000496360845312i^{10}i_2i_3 - 2891463543570311890566865090560i^{10}i_2i_3 +$   
 $88308178222163456i^{10}i_2i_3 + 599598052733571847066939392i^{10}i_2i_3 - 8811853774848i^{10}i_2i_3 -$   
 $14000309855379221647300i^{10}i_2i_3 + 97479251470446402968352866672640i^{10}i_2i_3$   
 $62914560i^{10}i_2i_3 + 3314614844289427456i^{10}i_2i_3 + 19743632405446596697071747072i^{10}i_2i_3 +$   
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 $2349627765367941191172096i^{10}i_2i_3 + 160080105674658429821334465478656i^{10}i_2i_3 +$   
 $4219485009161078228590656i^{10}i_2i_3 - 4663920783516479353440i^{10}i_2i_3 - 8349593486170324608i^{10}i_2i_3 -$   
 $25949272204749983774597552i^{10}i_2i_3 + 168462527977184i^{10}i_2i_3 + 3615767749123322594310432i^{10}i_2i_3$   
 $300837888i^{10}i_2i_3 - 215921145233199135744i^{10}i_2i_3 + 616229927403745317212702913984i^{10}i_2i_3 -$   
 $6144i^{10}i_2i_3 + 1066567715780800i^{10}i_2i_3 - 88753927451605190224113792i^{10}i_2i_3 +$

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 $859222867968i^{10}i_2i_3 - 158918678997153812471808i^{10}i_2i_3 + 603409881311665174411960989696i^{10}i_2i_3$   
 $8219687635411560887978987415249552000i^{10}i_2i_3 + 262144i^{10}i_2i_3 - 178652150455042048i^{10}i_2i_3 +$   
 $13743543044455536139034624i^{10}i_2i_3 + 4460718370970074913582229123077120i^{10}i_2i_3 +$   
 $755353629376i^{10}i_2i_3 - 31886510373696394197296i^{10}i_2i_3 - 5938917592170743659143396997632i^{10}i_2i_3 +$   
 $449342077792275016813671066186490000i^{10}i_2i_3 + 2584698285924232982401063716i^{10}i_2i_3 -$   
 $1076448164268941617256844i^{10}i_2i_3 + 114715810298410359519692i^{10}i_2i_3 - 3641229279555183859009414$   
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 $1396299980777369520070152i^{10}i_2i_3 + 10083133967459042825973517961268i^{10}i_2i_3 -$   
 $83047680i^{10}i_2i_3 - 3420729729797762816i^{10}i_2i_3 - 2328837190432645290652458416i^{10}i_2i_3 +$   
 $588794897426110537745808712861488i^{10}i_2i_3 - 256i^{10}i_2i_3 - 2818796161673088i^{10}i_2i_3 +$   
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 $79735296829i^{10}i_2i_3 + 18086829903613574584i^{10}i_2i_3 - 148503230450519874373700144i^{10}i_2i_3 -$   
 $412301901384828985965053850514920020i^{10}i_2i_3 - 2457676i^{10}i_2i_3 + 169510614703437600i^{10}i_2i_3 +$   
 $22795878549912015010928i^{10}i_2i_3 - 217552659984742507306371238199044i^{10}i_2i_3 +$   
 $6764692992i^{10}i_2i_3 + 434054607124859380768i^{10}i_2i_3 + 147786239372935408893012596736i^{10}i_2i_3 -$   
 $533359017011233109928i^{10}i_2i_3 - 1966085 + 1418294255707668448i^{10}i_2i_3 -$   
 $6324413850721664274452160i^{10}i_2i_3 + 14980184609949798113492863030784i^{10}i_2i_3 -$   
 $3883782492464615890560656006538925150000i^{10}i_2i_3$

# Computing modular polynomials

- 1 Dimension 1 : elliptic curves
- 2 Dimension 2 : abelian surfaces
  - Computation of the modular polynomials
  - Smaller invariants
- 3 Real Multiplication : cyclic isogenies

## Alternative invariants

⇒ look at modular functions for another group.

$$b_i(\Omega) := \frac{\theta_i(\Omega/2)}{\theta_0(\Omega/2)}, \quad i = 1, 2, 3.$$

Modular functions for  $\Gamma(2, 4)$ .

## Modular polynomials with $b_1, b_2, b_3$

### Theorem (Mumford, Manni)

*The field of modular functions invariant by  $\Gamma(2, 4)$  is  $\mathbb{C}(b_1, b_2, b_3)$ .*

We look at  $C_p = \Gamma(2, 4)/(\Gamma_0(p) \cap \Gamma(2, 4))$ ,  $p > 2$ .

The index is still  $p^3 + p^2 + p + 1$ .

### Proposition

*The field of modular functions invariant by  $\Gamma_0(p) \cap \Gamma(2, 4)$  is  $\mathbb{C}(b_1, b_2, b_3, b_{1,p})$ .*

We compute  $\Phi_{1,p}(X, b_1, b_2, b_3) = \prod_{\gamma \in C_p} (X - b_{1,p}^\gamma)$  and  $\Psi_{\ell,p}(X, b_1, b_2, b_3) = \sum_{\gamma \in C_p} b_{\ell,p}^\gamma \prod_{\gamma' \in C_p \setminus \{\gamma\}} (X - b_{1,p}^{\gamma'})$ . They are in  $\mathbb{Q}(b_1, b_2, b_3)[X]$ .

# Algorithm

- Evaluation of the  $b_i$  in  $\tilde{O}(N)$  (Dupont 2006, Enge–Thomé 2014).
- Inversion :  $(b_1, b_2, b_3)(\Omega) \longrightarrow \Omega ?$

$$(b_1, b_2, b_3)(\Omega) \longrightarrow (j_1, j_2, j_3)(\Omega) \longrightarrow \Omega \bmod \mathrm{Sp}_4(\mathbb{Z}).$$

**Problem : we want  $\Omega \bmod \Gamma(2, 4)$  !**

Solutions :

- Compute  $b_i(\gamma\Omega)$  for  $\gamma \in \mathrm{Sp}_4(\mathbb{Z})/\Gamma(2, 4)$ . But index 11520 !
- Use of functional equation of the theta functions.

## Denominators with the theta functions

Polynomials computed for  $p = 3, 5, 7$ .

Always  $D_p$  in the denominator.

$$\begin{aligned} D_3(b_1, b_2, b_3) = & 64(b_1^2 b_2^2 b_3^2)(16b_1^4 b_2^4 b_3^4 + 1)(b_1^4 + b_2^4 + b_3^4) \\ & - 32(48b_1^4 b_2^4 b_3^4 + 16b_1^2 b_2^2 b_3^2 + 1)(b_1^4 b_2^4 + b_1^4 b_3^4 + b_2^4 b_3^4) + \\ & 256(b_1^8 b_2^8 + b_1^8 b_3^8 + b_2^8 b_3^8) + 32(b_1^4 b_2^4 b_3^4)(-24b_1^4 b_2^4 b_3^4 + 80b_1^2 b_2^2 b_3^2 + 13) + 1. \end{aligned}$$

It is symmetric and there are relations modulo 2 and 4 between the exponents.

# Symmetries

## Theorem (M. 2014)

For all prime  $p$ , we have  $\Phi_{1,p}(X, b_1, b_2, b_3) = \Phi_{1,p}(X, b_1, b_3, b_2)$  and  $\Psi_{2,p}(X, b_1, b_3, b_2) = \Psi_{3,p}(X, b_1, b_2, b_3)$ .

**Proof** : there always exist  $\gamma \in \mathrm{Sp}_4(\mathbb{Z})/\Gamma(2, 4)$  such that for all  $\Omega \in \mathcal{H}_2$  :

$$\begin{array}{ll} b_1(\gamma\Omega) = b_1(\Omega) & b_{1,p}(\gamma\Omega) = b_{1,p}(\Omega) \\ b_2(\gamma\Omega) = b_3(\Omega) & \text{and} \quad b_{2,p}(\gamma\Omega) = b_{3,p}(\Omega) \\ b_3(\gamma\Omega) = b_2(\Omega) & b_{3,p}(\gamma\Omega) = b_{2,p}(\Omega) \end{array}$$

$\Phi_{1,p}$  is a minimal polynomial.

$\Psi_{\ell,p}(b_{1,p}) = b_{\ell,p} \Phi'_{1,p}(b_{1,p})$  for  $\ell = 2, 3$ . Action on  $\Psi_{2,p}(X)$  :

$$\Psi_{2,p}(b_{1,p}, b_1, b_3, b_2) = b_{3,p} \Phi'_{1,p}(b_{1,p}, b_1, b_2, b_3) := \Psi_{3,p}(b_{1,p}, b_1, b_2, b_3).$$



## Relations modulo 2 and 4

We look at matrices  $\gamma$  such that

$$\begin{array}{lcl} b_1(\gamma\Omega) & = & i^{\alpha_1} b_1(\Omega) \\ b_2(\gamma\Omega) & = & i^{\alpha_2} b_2(\Omega) \\ b_3(\gamma\Omega) & = & i^{\alpha_3} b_3(\Omega) \end{array} \quad \text{and} \quad \begin{array}{lcl} b_{1,p}(\gamma\Omega) & = & i^{\alpha_4} b_{1,p}(\Omega) \\ b_{2,p}(\gamma\Omega) & = & i^{\alpha_5} b_{2,p}(\Omega) \\ b_{3,p}(\gamma\Omega) & = & i^{\alpha_6} b_{3,p}(\Omega) \end{array}$$

## Comparison

For  $p = 3$  :

| $i$      | $j_1$ | $i_1$    | $b_1$ | $j_2$ | $i_2$    | $b_2$ | $j_3$ | $i_3$    | $b_3$ |
|----------|-------|----------|-------|-------|----------|-------|-------|----------|-------|
| 0        | 394   | 61       | 40    | 288   | 32       | 10    | 278   | 32       | 10    |
| 1        | 302   | 61       | 37    | 286   | 32       | 12    | 276   | 31       | 12    |
| 2        | 302   | 61       | 38    | 286   | 32       | 14    | 276   | 31       | 14    |
| $\vdots$ |       | $\vdots$ |       |       | $\vdots$ |       |       | $\vdots$ |       |
| 37       | 268   | 41       | 17    | 382   | 22       | 16    | 253   | 21       | 16    |
| 38       | 263   | 36       | 14    | 375   | 21       | 14    | 248   | 19       | 14    |
| 39       | 257   | 31       | 13    | 367   | 20       | 12    | 243   | 17       | 12    |

- $p = 3$  : 175 KB =  $\sim$  5000 smaller than Streng ;
- $p = 5$  : 200 MB ;
- $p = 7$  : 30 GB ;

# Computing modular polynomials

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# Hilbert space

Let  $D \in \mathbb{Z}^{>0}$  and  $K = \mathbb{Q}(\sqrt{D})$  a real quadratic field. We take  $D \in \{2, 5\}$  for simplicity.

The group  $\mathrm{SL}_2(\mathcal{O}_K)$  acts on  $\mathcal{H}_1^2$ .

## Proposition

The **Hilbert modular surface**  $\mathcal{H}_1^2/\mathrm{SL}_2(\mathcal{O}_K)$  is a moduli space for isomorphism classes of ppas with real multiplication by  $\mathcal{O}_K$ .

Let  $\rho$  a prime number such that

$$\rho = \beta\bar{\beta}, \quad \beta \in \mathcal{O}_K^+.$$

$\beta$ -isogenous surfaces :  $\beta\gamma z$ ,  $z \in \mathcal{H}_1^2$  and  $\gamma \in C_\rho = \mathrm{SL}_2(\mathcal{O}_K)/\Gamma_0(\beta)$ ;  
 $\#C_\rho = \rho + 1$ .

# Hilbert and Humbert

The following diagram is commutative :

$$\begin{array}{ccc} \mathcal{H}_1^2 & \xrightarrow{\phi} & \mathcal{H}_2 \\ \downarrow & & \downarrow \\ \mathcal{H}_1^2/\mathrm{SL}_2(\mathcal{O}_K) & \xrightarrow{\rho} & \mathcal{H}_2/\mathrm{Sp}_4(\mathbb{Z}) \end{array}$$

where  $\rho$  is generically of degree 2 onto the Humbert surface  $H_{\Delta_K}$ .

# Hilbert modular function

**Gundlach invariants** :  $J_1$  and  $J_2$  for  $D = 2$  and  $5$  only.

Two modular polynomials  $\Phi_\beta$  and  $\Psi_\beta$  in  $\mathbb{Q}(J_1, J_2)[X]$ .

# Algorithm

For  $D = 5$ , we have (Resnikoff 1974, Lauter–Yang 2011)

$$\begin{aligned}j_1 \circ \phi &= 8J_1(3J_2^2/J_1 - 2)^5; \\j_2 \circ \phi &= \frac{1}{2}J_1(3J_2^2/J_1 - 2)^3; \\j_3 \circ \phi &= 2^{-3}J_1(3J_2^2/J_1 - 2)^2(4J_2^2/J_1 + 2^5 3^2 J_2/J_1 - 3).\end{aligned}$$

These equations can be inverted by Gröbner basis.

Fast evaluation of the Gundlach invariants :

$$z \rightarrow \phi(z) = \Omega \rightarrow (j_1(\Omega), j_2(\Omega), j_3(\Omega)) \rightarrow (J_1(z), J_2(z)).$$

Inversion of the Gundlach invariants :

$$(J_1(z), J_2(z)) \rightarrow (j_1(\phi(z)), j_2(\phi(z)), j_3(\phi(z))) \rightarrow \phi(z) \rightarrow z.$$

# Results

**D=2**

|                     |       |       |       |      |      |       |
|---------------------|-------|-------|-------|------|------|-------|
| $p$                 | 2     | 7     | 17    | 23   | 31   | 41    |
| <i>Memory space</i> | 8.5KB | 172KB | 5.8MB | 21MB | 70MB | 225MB |

**D=5**

|                     |      |       |      |       |       |
|---------------------|------|-------|------|-------|-------|
| $p$                 | 5    | 11    | 19   | 29    | 31    |
| <i>Memory space</i> | 22KB | 3.5MB | 33MB | 188MB | 248MB |



# Theta functions

Other invariants?

$$\tilde{j}_i = j_i \circ \phi, \quad i = 1, 2, 3$$

or

$$\tilde{b}_i = b_i \circ \phi, \quad i = 1, 2, 3.$$

Works for any  $D$ . Three invariants for a space of dimension 2 : need the equation  $P$  of the Humbert component (Gruenewald 2008).

Interpolation :  $\mathbb{Q}(\tilde{b}_1, \tilde{b}_2, \tilde{b}_3)/(P) = \mathbb{Q}(\tilde{b}_1, \tilde{b}_2)[\tilde{b}_3]/(P)$ .

# Conclusion

- Implementation and generalization of the algorithm of Dupont ;
- Used smaller invariants and proved properties with them ;
- Definition and computation of modular polynomials with cyclic isogenies.

# Perspectives

- Compute more modular polynomials ;
- Release the code ;
- Applications of the polynomials.